



## Task Hiperkocka

*...it's dark in the cube, it's dark in the cube...*

Five in the morning. Daniel wakes up, he opens his eyes. His head hurts a bit. He can still hear the ringing in his ears.

He comes to realize that he has found himself at a playground, in a big metal box.

*...I was in the cube, I was in the cube...*

He remembers a similar situation he found himself in, three years ago, COCI round 2, task Kocka.

*...I'm in the cube again, I'm in the cube again...*

But this time, things are much more complicated... Daniel is in an  $n$ -dimensional hypercube  $Q_n$ .  $2^{n-1}$  identical copies of a tree  $\mathcal{T}$  with  $n$  edges are scattered around him. It soon became clear to him that salvation lies in tiling the edges of the hypercube with the trees.

Formally, a *hypercube*  $Q_n$  is a graph with nodes  $0, 1, \dots, 2^n - 1$ , in which nodes  $x$  and  $y$  are connected if and only if their bitwise *xor* is a power of two.

A tree can be *placed* on the hypercube so that:

- each node of the tree corresponds to some node of the hypercube
- those nodes have to be distinct
- if there is an edge between two nodes in the tree, then there has to be an edge between the corresponding nodes in the hypercube.

A *tiling* of the hypercube is done by placing several trees so that each edge of the hypercube belongs to at most one tree.

Your task is to tile the hypercube  $Q_n$  with as many copies of the given tree  $\mathcal{T}$ , which has  $n$  edges.

### Input

The first line contains a positive integer  $n$  ( $1 \leq n \leq 16$ ), the dimension of the hypercube.

Each of the following  $n$  lines contains two integers  $x$  and  $y$  ( $0 \leq x, y \leq n$ ,  $x \neq y$ ) which denote that the nodes  $x$  and  $y$  are connected by an edge in tree  $\mathcal{T}$ .

### Output

In the first line print the number of trees in your tiling.

Each of the following lines should describe a placement of a single copy of the tree  $\mathcal{T}$ .

In the  $i$ -th line print  $n + 1$  numbers  $a_0^{(i)}, a_1^{(i)} \dots a_n^{(i)}$ . These numbers denote that the  $i$ -th tree is placed so that the hypercube node  $a_j^{(i)}$  corresponds to the tree node  $j$ , for all  $j = 0, \dots, n$ .



## Scoring

If your solution correctly places  $k$  trees, you will receive  $f(k) \cdot 110$  points for that test case, where

$$f(k) = \begin{cases} 0.7 \cdot k/2^{n-1} & \text{if } k < 2^{n-1} \\ 1 & \text{if } k = 2^{n-1}. \end{cases}$$

Of course, if your solution is not correct, you will receive 0 points.

Your total number of points is equal to the minimum number of points your solution receives over all of the test cases.

It is possible to prove that there always exists a solution which uses all of the  $2^{n-1}$  trees.

## Examples

**input**

1  
0 1

**output**

1  
0 1

**input**

2  
0 1  
1 2

**output**

2  
0 1 3  
0 2 3

**input**

3  
0 1  
0 2  
0 3

**output**

4  
0 1 2 4  
3 1 2 7  
5 1 4 7  
6 2 4 7

Clarification of the third example:

